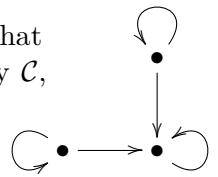


Minicourse on Category Theory: Exercise sheet 1

Exercise 1.1. *First steps with categories*

- (a) Write down the category for your favorite subject of mathematics or computer science: What are its objects, what are its morphisms? What are its initial and terminal objects?
- (b) Write down a formal definition of the category on the right: What precisely are the objects and the morphisms? What is the composition rule?
- (c) Prove: Identity morphisms in categories are unique. More precisely, verify that if id_X and $\tilde{\text{id}}_X$ are both identity morphisms for some object X in a category \mathcal{C} , then $\text{id}_X = \tilde{\text{id}}_X$.



Exercise 1.2. *Initial and terminal objects*

Determine the initial and terminal objects (in case they exist) of ...

- (a) the category $\text{Vect}(\mathbb{R})$
- (b) the numerical category
- (c) the category of Pokémon
- (d) the category of Exercise 1.1(b)

Exercise 1.3. *Groups as categories*

Explain how out of any group G (or even just a monoid) a category can be constructed, in such a way that the composition rule of G plays a role.

Remark. The resulting category is sometimes denoted “ BG ” and arises in algebraic topology.

Exercise 1.4. *Uniqueness up to unique isomorphism*

An *isomorphism* $f : X \rightarrow Y$ in a category is a morphism for which a morphism $g : Y \rightarrow X$ such that $f \circ g = \text{id}_Y$ and $g \circ f = \text{id}_X$ exists. We also write “ f^{-1} ” for g .

- (a) Prove: The isomorphisms in the category of sets are precisely the bijective maps.
- (b) List two different terminal objects of Set .
- (c) Let T and T' be terminal objects of a category \mathcal{C} . Show that there is an isomorphism between T and T' , and moreover that there is exactly one such isomorphism.

Minicourse on Category Theory: Exercise sheet 2

Definition. A *monomorphism* in a category \mathcal{C} is a morphism $f : X \rightarrow Y$ such that for all objects Z and all morphisms $p, q : Z \rightarrow X$,
if $f \circ p = f \circ q$, then $p = q$.

Dually, an *epimorphism* is a morphism $f : X \rightarrow Y$ such that for all objects Z and all morphisms $p, q : Y \rightarrow Z$, if $p \circ f = q \circ f$, then $p = q$.

Exercise 2.1. Monomorphisms and epimorphisms in the category of sets

Let $f : X \rightarrow Y$ be a map between sets. Prove:

- (a) The map f is a monomorphism in \mathbf{Set} if and only if it is injective.
- (b) The map f is an epimorphism in \mathbf{Set} if and only if it is surjective.

Remark. This exercise is surprisingly deep from the point of view of mathematical logic: Can you give a direct proof avoiding the law of excluded middle, a proof which could be formalized in \mathbf{IZF} ? Can you even give a proof avoiding the powerset axiom?

Exercise 2.2. Composition of monomorphisms

Let \mathcal{C} be a category.

- (a) Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ be morphisms in \mathcal{C} . Show: If $g \circ f$ is monic, so is f .
- (b) Your proof of (a) works in all categories, hence also in \mathcal{C}^{op} . What does it state then?

Exercise 2.3. Balanced categories

- (a) Prove: The isomorphisms in the category of monoids are precisely what is commonly known as monoid isomorphisms (bijective monoid homomorphisms).
- (b) Prove: Isomorphisms are monic and epic.

Remark. Categories in which the converse holds are called *balanced*. Examples for balanced categories are \mathbf{Set} , \mathbf{Grp} , $\mathbf{Vect}(\mathbb{R})$ and many others; however there are also many categories which are not balanced.

- (c) Give an example for a category with morphisms which are monic and epic, but fail to be isomorphisms.

Exercise 2.4. The “2 out of 3” property

- (a) Let f be a morphism in a category which has both a left and a right inverse. Prove that f is an isomorphism and that the given one-sided inverses are equal.

Remark. A *left inverse* to $f : X \rightarrow Y$ is a morphism $g : Y \rightarrow X$ such that $g \circ f = \text{id}_X$. Dually, a *right inverse* to f is a morphism $h : Y \rightarrow X$ such that $f \circ h = \text{id}_Y$.

- (b) Let $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ be morphisms in a category. Prove that, if two out of the three morphisms f , g and $g \circ f$ are isomorphisms, so is the third.

Minicourse on Category Theory: Exercise sheet 3

Exercise 3.1. *Product with the terminal object*

Let X be an object of a category \mathcal{C} .

- (a) Assume that \mathcal{C} contains a terminal object 1 . Prove:

$$X \times 1 \cong X.$$

However, this claim is not to be taken literally. Instead, verify more precisely that X (together with which morphisms?) serves as a product of X and 1 .

- (b) What is the dual statement to (a)?

Exercise 3.2. *Category of wannabe products*

Let X and Y be objects of a category \mathcal{C} . We define the following category of *wannabe products* of X and Y :

objects: diagrams of the form $X \leftarrow Q \rightarrow Y$ in \mathcal{C}
morphisms: $\text{Hom}(X \leftarrow Q \rightarrow Y, X \leftarrow R \rightarrow Y) :=$

$$\left\{ \begin{array}{c} \text{commutative diagrams of the form} \end{array} \begin{array}{ccccc} & & Q & & \\ & \swarrow & \downarrow & \searrow & \\ X & & & & Y \\ & \nwarrow & \downarrow & \nearrow & \\ & & R & & \end{array} \right\}$$

- (a) How should the composition rule be defined in this diagram category?
(b) Explain that a product of X and Y in \mathcal{C} is the same as a terminal object in the category of wannabe products of X and Y .
(c) Which corollary follows from (b) in view of Exercise 1.4?

Exercise 3.3. *Infima in preorders*

Let X be a *preorder*, that is a set X together with a reflexive and transitive (but not necessarily antisymmetric) binary relation (\preceq) on X . For instance \mathbb{Z} with the divisibility ordering is a preorder.

- (a) Construct a sensible category BX out of X . Why are the category axioms satisfied?
(b) When are two objects of BX isomorphic?
(c) An *infimum* of two elements $a, b \in X$ is an element $p \in X$ such that

$$\forall x \in X. \quad x \preceq a \text{ and } x \preceq b \iff x \preceq p.$$

Show: An infimum of a and b is the same as a product of a and b in BX .

Minicourse on Category Theory: Exercise sheet 4

Exercise 4.1. *Functors as generalization of group homomorphisms*

- (a) Explain how a group homomorphism $\varphi : G \rightarrow H$ gives rise to a functor $B(\varphi) : BG \rightarrow BH$.

Remark. For a group X , the category BX is the category which has a single object \heartsuit and one morphism $g : \heartsuit \rightarrow \heartsuit$ for each group element $g \in X$, with id_{\heartsuit} being the neutral element of X and composition given by the operation of the group.

- (b) Upgrade your solution of part (a) to a functor $\text{Grp} \rightarrow \text{Cat}$.

Exercise 4.2. *Functors between categories induced by preorders*

Let X and Y be preorders. Let BX and BY be the induced thin categories of Exercise 3.3. By which data is a functor $BX \rightarrow BY$ given? State and prove an analogue to Exercise 4.1.

Exercise 4.3. *Properties of functors*

- (a) Study three functors of your choosing as to whether they are full, faithful or essentially surjective.
- (b) Prove that the numerical category is equivalent to its opposite category.

Remark. The numerical category has as objects the natural numbers, and the set of morphisms from m to n is by definition the set of real $(n \times m)$ -matrices.

- (c) Prove that the functor from the numerical category to the category of finite-dimensional real vector spaces, mapping n to \mathbb{R}^n and a matrix to its induced linear transformation, is fully faithful and essentially surjective.

Exercise 4.4. *Functoriality of taking products*

Let A be an object of a category \mathcal{C} . Assume that we are given, for each object X of \mathcal{C} , a product $A \times X$ of A and X . Explain how the incomplete association

$$X \longmapsto A \times X$$

can be extended to a well-defined functor $F : \mathcal{C} \rightarrow \mathcal{C}$: How should F be defined on morphisms? Why are the functor axioms satisfied?

Exercise 4.5. *In preparation of the Yoneda lemma I*

Let A be an object of a (locally small) category \mathcal{C} . The *contravariant Hom-functor associated to A* is defined by

$$\begin{aligned} \widehat{A} : \quad \mathcal{C}^{\text{op}} &\longrightarrow \text{Set} \\ X &\longmapsto \text{Hom}_{\mathcal{C}}(X, A) \\ (f : X \rightarrow Y) &\longmapsto f^*, \end{aligned}$$

where f^* is the map

$$\begin{aligned} f^* : \text{Hom}_{\mathcal{C}}(X, A) &\longrightarrow \text{Hom}_{\mathcal{C}}(Y, A) \\ g &\longmapsto g \circ f. \end{aligned}$$

Verify that \widehat{A} is indeed a functor.

Minicourse on Category Theory: Exercise sheet 5

Exercise 5.1. *Non-equivalent categories*

- (a) Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be an elementary equivalence of categories (fully faithful, essentially surjective). Show that an object X of \mathcal{C} is initial if and only if $F(X)$ is initial in \mathcal{D} .
- (b) Find categorical properties which distinguish all of the following categories:

$$\mathbf{Set}, \quad \mathbf{Vect}(\mathbb{R}), \quad \mathbf{Grp}, \quad \mathbf{Pok}, \quad \mathbf{Set}^{\text{op}}.$$

Remark. The category \mathbf{Set}^{op} obtains a more concrete description by realizing that it is equivalent, by the contravariant powerset functor, to the category of *complete atomic Heyting algebras*. The (only) examples of such algebras are powersets. However this remark is not particularly relevant to the exercise.

Exercise 5.2. *Length as a natural transformation*

Let L be the “underlying functor of the list monad”, that is the functor $L : \mathbf{Set} \rightarrow \mathbf{Set}$ which maps a set X to the set $L(X)$ of finite lists of elements of X ,

$$L(X) = \{[x_1, x_2, \dots, x_n] \mid n \in \mathbb{N}, x_1, \dots, x_n \in X\},$$

and a map f to the induced map $L(f)$ given by $L(f)([x_1, \dots, x_n]) = [f(x_1), \dots, f(x_n)]$. For each set X , let $\eta_X : L(X) \rightarrow \mathbb{N}$, $\mathbf{xs} \mapsto (\text{length of } \mathbf{xs})$.

- (a) Show that the maps $(\eta_X)_X$ form a natural transformation $L \Rightarrow K_{\mathbb{N}}$, where $K_{\mathbb{N}}$ is the constant functor on \mathbb{N} .
- (b) Find a natural transformation $L \circ L \Rightarrow L$ and a natural transformation $\text{Id}_{\mathbf{Set}} \Rightarrow L$.

Remark. Recall that $\text{Id}_{\mathbf{Set}} : \mathbf{Set} \rightarrow \mathbf{Set}$ is the identity functor on \mathbf{Set} , given by $X \mapsto X$ and $f \mapsto f$.

- (c) What would go wrong with finite sets instead of finite lists?

Exercise 5.3. *Examples for natural transformations*

- (a) Show: There is just one natural transformation $\eta : \text{Id}_{\mathbf{Set}} \Rightarrow \text{Id}_{\mathbf{Set}}$, namely the family consisting of the maps $\eta_X : X \rightarrow X$, $x \mapsto x$.

Hint (also for (b)). Consider suitable maps $\{\heartsuit\} \rightarrow X$.

- (b) Let $D : \mathbf{Set} \rightarrow \mathbf{Set}$ be the functor $X \mapsto X \times X$, $f \mapsto f \times f := ((a, b) \mapsto (f(a), f(b)))$. Show: There is just one natural transformation $\omega : \text{Id}_{\mathbf{Set}} \Rightarrow D$, namely

$$\omega_X : X \rightarrow X \times X, \quad x \mapsto (x, x).$$

- (c) Let $P : \mathbf{Set} \rightarrow \mathbf{Set}$ be the (covariant) powerset functor. Show: There is no natural transformation $P \Rightarrow \text{Id}_{\mathbf{Set}}$, but a natural transformation in the other direction.
- (d) Assume that we are given, for each inhabited set X , an element $a_X \in X$. Show: The maps $\tau_X : X \rightarrow X$, $x \mapsto a_X$ do *not* constitute a natural transformation $\text{Id}_{\mathcal{C}} \Rightarrow \text{Id}_{\mathcal{C}}$, where \mathcal{C} is the category of inhabited sets and arbitrary maps.
- (e) Which natural transformations $\text{Id}_{\mathbf{Vect}(\mathbb{R})} \Rightarrow \text{Id}_{\mathbf{Vect}(\mathbb{R})}$ are there?

Exercise 5.4. *In preparation of the Yoneda lemma II*

Let $\varphi : A \rightarrow B$ be a morphism in a (locally small) category \mathcal{C} . Construct a natural transformation $\hat{A} \Rightarrow \hat{B}$ between the Hom functors of Exercise 4.5.

Minicourse on Category Theory: Exercise sheet 6

Exercise 6.1. *Natural transformations between monotone maps*

Let $f, g : X \rightarrow Y$ be monotone maps between preorders X and Y . Let Bf and Bg be the induced functors between the thin categories BX and BY as in Exercise 4.2. Characterize when there is a natural transformation $Bf \Rightarrow Bg$.

Exercise 6.2. *Pseudoinverses of elementary equivalences*

Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be an elementary equivalence, so a fully faithful essentially surjective functor. Assume that for each object $Y \in \mathcal{D}$, an object $X_Y \in \mathcal{C}$ with $F(X_Y) \cong Y$ is given. Explain how the mapping $Y \mapsto X_Y$ can be extended to a functor $G : \mathcal{D} \rightarrow \mathcal{C}$, and show that $G \circ F \cong \text{Id}_{\mathcal{C}}$ and $F \circ G \cong \text{Id}_{\mathcal{D}}$.

Exercise 6.3. *Linear orderings and permutations*

For a finite set X , let $\text{Perm}(X)$ be the set of bijections from X to X and let $\text{LinOrd}(X)$ be the set of linear orderings on X (lists $[x_1, \dots, x_n]$ in which every element of X occurs exactly once). Let \mathcal{S} be the category of *species*, that is the subcategory of Set which has as objects the finite sets and as morphisms the bijections.

- (a) Show that for every finite set X , $\text{Perm}(X) \cong \text{LinOrd}(X)$.
- (b) Explain how Perm and LinOrd can be made into functors $\mathcal{S} \rightarrow \text{Set}$.
- (c) Show that there is no natural transformation $\text{Perm} \Rightarrow \text{LinOrd}$ and hence also no natural isomorphism $\text{Perm} \cong \text{LinOrd}$.

Hint. Consider the identity permutation.

Exercise 6.4. *Tip already in the diagram*

Let $F : \mathcal{I} \rightarrow \mathcal{C}$ be an \mathcal{I} -shaped diagram in a category \mathcal{C} . Assume that \mathcal{I} contains a terminal object T . Prove that F has a colimit, namely a certain cocone with tip $F(T)$.

Exercise 6.5. *Polynomials and power series*

Let K be a field (or just a commutative ring) and let $K[X]_n$ be the vector space of polynomials of degree $\leq n$ with coefficients in K .

- (a) Verify that $K[X]$ can be made into a colimit of the diagram $(K[X]_0 \hookrightarrow K[X]_1 \hookrightarrow K[X]_2 \hookrightarrow \dots)$, where the morphisms are all the inclusions.
- (b) Verify that the vector space $K[[X]] := \{\sum_{i=0}^{\infty} a_i X^i \mid a_0, a_1, \dots \in K\}$ of formal power series can be made into a limit of the diagram $(\dots \twoheadrightarrow K[X]_2 \twoheadrightarrow K[X]_1 \twoheadrightarrow K[X]_0)$, where the maps truncate the highest coefficients.

Minicourse on Category Theory: Exercise sheet 7

$$\begin{array}{ccc} X & \xrightarrow{\text{id}} & X \\ \text{id} \downarrow & \lrcorner & \downarrow f \\ X & \xrightarrow{f} & Y \end{array}$$

Exercise 7.1. Monomorphisms as limits

- (a) Verify that a morphism $f : X \rightarrow Y$ in a category is monic if and only if the diagram on the right is a pullback square.
- (b) Let $\eta : F \Rightarrow G$ be a natural transformation between functors $F, G : \mathcal{C} \rightarrow \mathcal{D}$. Assume that \mathcal{D} has all fiber products. Prove that η is a monomorphism in $[\mathcal{C}, \mathcal{D}]$ if and only if all components η_X are monomorphisms in \mathcal{D} .

Hint. Limits in functor categories are computed objectwise.

Exercise 7.2. The Yoneda lemma

Let \mathcal{C} be a (small) category and let $\text{PSh}(\mathcal{C}) := [\mathcal{C}^{\text{op}}, \text{Set}]$ be its presheaf category. Our goal is to prove the *Yoneda lemma*, stating that there is a bijection

$$\text{Hom}_{\text{PSh}(\mathcal{C})}(\hat{X}, F) \cong F(X), \quad (\star)$$

natural in $X \in \text{Ob}(\mathcal{C})$ and $F \in \text{Ob}(\text{PSh}(\mathcal{C}))$, where \hat{X} is the functor from Exercise 4.5.

- (a) Prove that a natural transformation $\eta : \hat{X} \Rightarrow F$ is already uniquely determined by the single value $s := \eta_X(\text{id}_X) \in F(X)$, namely by the relationship

$$\eta_Y(f) = F(f)(s) \quad (\dagger)$$

for all objects Y and morphisms $f \in \text{Hom}_{\mathcal{C}}(Y, X)$.

- (b) Show that conversely, every $s \in F(X)$ defines a natural transf. $\eta : \hat{X} \Rightarrow F$ by (\dagger) .
- (c) Verify that there is a bijection (\star) for fixed $X \in \text{Ob} \mathcal{C}$ and $F \in \text{Ob} \hat{\mathcal{C}}$.
- (d) Domain and codomain of (\star) can be regarded as evaluations of the functors

$$\begin{aligned} L : \mathcal{C}^{\text{op}} \times \text{PSh}(\mathcal{C}) &\longrightarrow \text{Set}, & (X, F) &\longmapsto \text{Hom}_{\text{PSh}(\mathcal{C})}(\hat{X}, F) \\ R : \mathcal{C}^{\text{op}} \times \text{PSh}(\mathcal{C}) &\longrightarrow \text{Set}, & (X, F) &\longmapsto F(X) \end{aligned}$$

at (X, F) . Explain how these functors act on morphisms, and prove that they are isomorphic, concluding your proof of Yoneda's lemma. Congratulations!